

Shuffle-Raitonal Series, Doubly Rational Series & Shuffle-Finite Series

Subbarao Venkatesh Guggilam

Department of Mathematics, Informatics and Mechanics
University of Warsaw, Poland.

SAMSA, 2026

- 1 A brief of Chen–Fließ Series
- 2 Shuffle Algebra
- 3 Shuffle-Rational Series
- 4 Action of Rational Functions on Formal Series

Notations

- $X = \{x_0, x_1, x_2, \dots, x_m\}$ be a finite set of noncommuting alphabet.
- X^* is the free monoid generated by X . The elements of X are words, with empty word denoted by \emptyset . The subset of nonempty words viz. $X^+ := X^* \setminus \{\emptyset\}$.
- $\mathbb{R}\langle\langle X \rangle\rangle$ is the linear dual to the free algebra generated by X . It is the (Cauchy) algebra of all formal series, where $\mathbb{R}\langle\langle X \rangle\rangle \ni c : X^* \rightarrow \mathbb{R}$
- A typical way to represent the formal series c is

$$c = \sum_{w \in X^*} c(w)w$$

- $\mathbb{R}\langle X \rangle$: noncommutative polynomials is the subalgebra of $\mathbb{R}\langle\langle X \rangle\rangle$ that have finite support.

What is a Chen–Fliess Series ?

Every series $c \in \mathbb{R}\langle\langle X \rangle\rangle$ has an associated Chen–Fliess series

$$F_c[u](t) = \sum_{\eta \in X^*} c(\eta) F_\eta[u](t), \quad (1)$$

where $E_\emptyset[u] = 1$ and

$$F_{x_i \eta}[u](t) = \int_0^t d\tau u_i(\tau) F_\eta[u](\tau)$$

with $x_i \in X$, $\eta \in X^*$, $u \in L_1^m[0, t]$ and $u_0 = 1$ (Fliess, 1981).

Why should anyone care about Chen–Fliess Series?

“Nonlinear dynamical systems” which are modeled by controlled differential equations on \mathbb{R}^n is given by

$$\dot{z} = f(z) + \sum_{i=1}^m g_i(z)u_i(t), \quad z(0) = z_0 \in \mathbb{R}^n$$

$$y(t) = h(z(t))$$

with real analytic vector fields f, g_1, g_2, \dots, g_m and with a real analytic function h (around z_0),

The local input-output behaviour of such systems viz. a map from $\{u_1, u_2, \dots, u_m\} \mapsto y$ is modeled by a Chen–Fliess series. (Fliess, 1981)

A small goodie bag!

One can ask for the converse:

What conditions on the formal power series $c \in \mathbb{R}\langle\langle X \rangle\rangle$ are necessary and sufficient that the corresponding Chen–Fliess series has a canonical realization as a local input-to-output map of a controlled nonlinear system ?

Fliess & Reutenauer (1982) answered this question:

- The coefficients of c has a Cauchy Growth
- Lie rank of c is finite,

Shuffle Product

The shuffle product of two words is inductively defined as

$$(x_i\eta) \sqcup (x_j\xi) = x_i(\eta \sqcup (x_j\xi)) + x_j((x_i\eta) \sqcup \xi),$$

where $\eta, \xi \in X^*$, $x_i, x_j \in X$, and $\eta \sqcup \emptyset = \emptyset \sqcup \eta = \eta$.

Example

Examples

$$x_1 \sqcup x_2 = x_1x_2 + x_2x_1$$

$$x_1x_2 \sqcup x_1 = 2x_1^2x_2 + x_1x_2x_1$$

$$x_1 \sqcup x_1 = 2x_1^2$$

Observe that the shuffle product is commutative by definition.

The shuffle product is then extended bilinearly to series in $\mathbb{R}\langle\langle X \rangle\rangle$. Say for $c, d \in \mathbb{R}\langle\langle X \rangle\rangle$

$$(c \sqcup d)(\eta) = \sum_{\substack{u, v \in X^* \\ \eta \in \text{supp}(u \sqcup v)}} c(u)d(v)$$

$(\mathbb{R}\langle\langle X \rangle\rangle, \sqcup)$ is an associative and commutative algebra.

For $c, d \in \mathbb{R}\langle\langle X \rangle\rangle$, the product of their corresponding Chen–Fliess series is a Chen–Fliess series given by the shuffle product of c and d (Ree, 1958)

$$F_c \cdot F_d[u] = F_{c \sqcup d}[u]$$

Shuffle Group

The subset of non proper series (constant term nonzero) in $\mathbb{R}\langle\langle X \rangle\rangle$ constitutes a group under the shuffle product.

The shuffle inverse of a non proper series is taken to be

$$c^{\sqcup^{-1}} = (c(\emptyset)(1 - c'))^{\sqcup^{-1}} = c(\emptyset)^{-1}(c')^{\sqcup^*}, \quad (2)$$

where $c' := 1 - c/c(\emptyset)$ is proper, and $(c')^{\sqcup^*} = \sum_{k \in \mathbb{N}_0} (c')^{\sqcup k}$. Here

$(c')^{\sqcup k} \triangleq c' \sqcup (c')^{\sqcup k-1}$ with $(c')^{\sqcup 0} = 1\emptyset$.

Example

Let $c = 1 - x_1 \in \mathbb{R}\langle\langle X \rangle\rangle$ so that $c' = x_1$. Then

$$c^{\sqcup^{-1}} = x_1^{\sqcup^*} = \sum_{k \in \mathbb{N}_0} k! x_1^k.$$

Definition

An \mathbb{R} -subalgebra \mathcal{F} of an \mathbb{R} -algebra on $\mathbb{R}\langle\langle X \rangle\rangle$ is said to be *rationally closed* if and only if the inverse of all invertible elements of \mathcal{F} belongs to \mathcal{F} .

The *rational closure* of an \mathbb{R} -subalgebra \mathcal{F}' of an \mathbb{R} -algebra on $\mathbb{R}\langle\langle X \rangle\rangle$ is the smallest rationally closed subalgebra \mathcal{F} containing \mathcal{F}' .

Classically, *rational series* are defined to be those in the rational closure of the \mathbb{R} -subalgebra of polynomials $\mathbb{R}\langle X \rangle$, where the \mathbb{R} -algebra structure on $\mathbb{R}\langle\langle X \rangle\rangle$ is under Cauchy product (Berstel & Reutenauer).

Rational Series

The following statements are known to be equivalent:

- Series c is rational.
- Series c is recognizable (there exists a linear representation λ, μ, γ) where $\lambda, \gamma \in \mathbb{R}^n$ and $\mu : X^* \rightarrow M_n(\mathbb{R})$ is a monoid morphism and $\forall w \in X^*$

$$c(w) = \lambda^t \mu(w) \gamma$$

(Schützenberger Theorem) (Schützenberger, 1961).

- Series c has a Hankel matrix with finite rank n (Fliess, 1974).
- Operator $y = F_c[u]$ has a canonical bilinear state space realization (Fliess, 1981)

$$\dot{z} = (\mu(x_0) + \sum_{i=1}^m \mu(x_i) u_i(t)) \quad z(0) = \gamma$$
$$y = \lambda^t z$$

All is well, but where is the question?

- What is the rational closure of subalgebra of polynomials in the shuffle algebra ?
- Can one give “decent” notion of recognizability ?
- What is the corresponding canonical realization as a nonlinear system ?
- What are the other possible equivalent notions ?
- What is the machine that recognizes rational closure of polynomials in shuffle algebra ?

Definition

A series c is said to be shuffle rational if it is in the rational closure of $\mathbb{R}\langle X \rangle$ in $(\mathbb{R}\langle\langle X \rangle\rangle, \sqcup)$.

Equivalently, c can be written as finite number of shuffle products of polynomials shuffled with finite number of shuffle inverse of polynomials

$$c = (c_1 \sqcup c_2 \sqcup \cdots c_k) \sqcup (d_1 \sqcup d_2 \sqcup \cdots d_m)^{\sqcup -1}$$

where $\{c_i, d_j\}_{\substack{j=1:k \\ j=1:m}} \in \mathbb{R}\langle X \rangle$.

Rational Functions on Formal Series

Let $p, q \in \mathbb{R}[y]$ and $c \in \mathbb{R}\langle\langle X \rangle\rangle$. Assume $p(y) = \sum_{i=0}^k a_i y^i$, where $k \in \mathbb{N}_0$:

The composition of p and c is defined as

$$p(c) = \sum_{i=0}^k a_i c^{\sqcup i}.$$

Extending the definition to rational functions $\frac{p}{q} \in \mathbb{R}(y)$ gives

$$\frac{p}{q}(c) = p(c) \sqcup q(c)^{\sqcup -1}.$$

provided $c(\emptyset)$ is not root of $q(y)$

These definitions can be generalized to action of multivariate polynomials on k -tuples of series as the shuffle product is \mathbb{R} -bilinear. Say

$$p = y_1 y_2 \in \mathbb{R}[y_1, y_2, \dots, y_k]$$

$$y_1 y_2(c_1, c_2, \dots, c_k) = c_1 \sqcup c_2$$

where and $c_i \in \mathbb{R}\langle\langle X \rangle\rangle$, $i = 1, 2, \dots, k$.

In general, given a rational function $\frac{p}{q} \in \mathbb{R}(y_1, y_2, \dots, y_n)$ and formal series $c_1, c_2, \dots, c_k \in \mathbb{R}\langle\langle X \rangle\rangle$, then

$$\frac{p}{q}(c_1, c_2, \dots, c_k) = p(c_1, c_2, \dots, c_k) \sqcup q(c_1, c_2, \dots, c_k)^{\sqcup -1}$$

provided $(c_1(\emptyset), c_2(\emptyset), \dots, c_k(\emptyset))$ is not a root of the denominator polynomial q .

Definition (Shuffle-Recognizable)

A series $c \in \mathbb{R}\langle\langle X \rangle\rangle$ is said to be *shuffle-recognizable* if

$c = p/q \left(\sum_{w \in X^*} \lambda^T \mu(w) \gamma w \right)$, where $\lambda^T = (\lambda_1^T \times \lambda_2^T \times \dots \times \lambda_k^T)$, $\mu = (\mu_1 \times \mu_2 \times \dots \times \mu_k)$, and $\gamma = (\gamma_1 \times \gamma_2 \times \dots \times \gamma_k)$. Here $\lambda_i, \gamma_i \in \mathbb{R}^{n_i}$ and $\mu_i : X^* \rightarrow N_n(\mathbb{R})$ is a monoid morphism into strictly upper triangular matrices.

The tuple $(p, q, \{\lambda_i\}_{i=1}^k, \{\mu_i\}_{i=1}^k, \{\gamma_i\}_{i=1}^k)$ is called a k^{th} -order *shuffle-representation* of c

Theorem (V. & Gray (2020))

Shuffle Rational = Shuffle Recognizable.

Theorem (V. & Gray (2020))

- A series c is shuffle-rational
- The Chen–Fliess series F_c has a (not necessarily canonical) realization as a nonlinear system

$$\dot{z} = (BD(\mu_1(x_0), \dots, \mu_k(x_0))) + \sum_{j=0}^m (BD(\mu_1(x_j), \dots, \mu_k(x_j)))u_j(t))$$

$$y = \frac{p}{q}(BD(\lambda_1, \dots, \lambda_k)^t z)$$

where $z(0) = \text{col}(\gamma_1, \dots, \gamma_k)$

References

- M. P. Schützenberger, On the definition of a family of automata, *Inf. Control*, 4 (1961) 245–270.
- M. Fliess, Matrices de Hankel, *J. Math. Pures Appl.*, 53 (1974) 197–224, +erratum, 54 (1975) 481.
- M. Fliess, Fonctionnelles causales non linéaires et indéterminées non commutatives, *Bull. Soc. Math. France*, 109, 1981, 3–40.
- M. Fliess and C. Reutenauer, Une application de l’algèbre différentielle aux systèmes réguliers (ou bilinéaires), in *Analysis and Optimization of Systems*, A. Bensoussan and J. L. Lions, Eds, Lect. Notes Control Informat. Sci., vol. 44, Springer-Verlag, Berlin, 1982, pp. 99–107.
- Y. Wang and E. D. Sontag, Algebraic differential equations and rational control systems, *SIAM J. Control Optim.*, 30 (1992) 1126–1149.
- G. S. Venkatesh and W. S. Gray, Shuffle-Rational Series: Recognizability and Realizations, *Proc. 24th Int. Conf. on System Theory, Control and Computing*, Sinaia, Romania, 2020, pp. 404–411.

Thank You
Questions?